

# SUGGESTED SOLUTION

SYJC

**SUBJECT- MATHS & STATS** 

Test Code – SYJ 6097

BRANCH - () (Date :)

Head Office : Shraddha, 3<sup>rd</sup> Floor, Near Chinai College, Andheri (E), Mumbai – 69. Tel : (022) 26836666

## ANSWER 1(A)

Given :  $l_2 = 80$ ,  $d_2 = 5$ ,  $q_3 = 0.20$ ,  $l_3 = ?$ ,  $l_4 = ?$ We have,  $d_x = l_x - l_{x+1}$   $\therefore d_2 = l_2 - l_3$   $\therefore 5 = 80 - l_3$   $\therefore l_3 = 80 - 5 = 75$ We have,  $q_x = \frac{d_x}{l_x}$   $\therefore q_3 = \frac{d_3}{l_3}$   $\therefore 0.20 = \frac{d_3}{75}$   $\therefore 0.20 \times 75 = d_3$   $\therefore d_3 = 15$ Now,  $d_3 = l_3 - l_4$   $\therefore 15 = 75 - l_4$   $\therefore l_4 = 75 - 15$  $\therefore l_4 = 60$ 

Hence,  $I_3 = 75$  and  $I_4 = 60$ 

## ANSWER 1(B)

Using the optimal sequence algorithm, we can obtain the following optimal sequence.

E B D C A

Next, we find total elapsed time as follows :

Job	Machine M <sub>1</sub>		Mach	ine M₂	Idle time for
	Time in	Time out	Time in	Time Out	M <sub>2</sub>
E	0	6	6	17	06
В	6	19	19	34	02
D	19	31	34	40	00
C	31	38	40	45	00
А	38	42	45	48	00
				Total	08

Minimum elapsed (Processing) time T = 48 hrs.

## ·.

(3 MARKS)

Idle time for machine  $M_1$ = T - (sum of the processing time of the five jobs on  $M_1$ ) = 48 - 42 = 06 hrs. Idle time for machine  $M_2$ = T - (sum of the processing time of the five jobs on  $M_2$ ) = 48 - 40 = 08 hrs.

(3 MARKS)

## ANSWER 2(A)

Let  $l_x$  = number of persons living at age x We know that,

$$d_{x} = l_{x} - l_{x+1} \qquad \dots (1)$$

$$p_{x} = \frac{l_{x+1}}{l_{x}} \text{ and also } p_{x} = \frac{e_{x}}{1 + e_{x+1}}$$

$$\therefore \frac{l_{x+1}}{l_{x}} = p_{x} = \frac{e_{x}}{1 + e_{x+1}}$$

$$\therefore l_{x+1} = \frac{e_{x} \cdot l_{x}}{1 + e_{x+1}}$$
Using this result in (1) we get
$$d_{x} = l_{x} = \frac{e_{x} \cdot l_{x}}{1 + e_{x+1}} \qquad \dots (2)$$
Given :  $d_{1} = 10, e_{1}^{0} = 3.22, e_{2}^{0} = 2.56$ 
We have  $e_{x}^{0} = e_{x} + \frac{1}{2}$ 

$$\therefore e_{x} = e_{x}^{0} - \frac{1}{2} = e_{x}^{0} - 0.5 \qquad \dots (3)$$

$$e_{1}^{0} = 3.22 \text{ and } e_{2}^{0} = 2.56$$
Using the result (3) we get
$$e_{1} = 3.22 - 0.5 = 2.72$$

$$e_{2} = 2.56 - 0.5 = 2.06$$
From (2),
$$d_{1} = l_{1} \left[ 1 - \frac{e_{1}}{1 + e_{2}} \right]$$

Putting  $d_1 = 10$ ,  $e_1 = 2.72$  and  $e_2 = 2.06$ , we get

$$10 = l_1 \left[ 1 - \frac{2.72}{1+2.06} \right]$$
  

$$\therefore 10 = l_1 \left[ 1 - \frac{2.72}{3.06} \right]$$
  

$$\therefore 10 = l_1 \left[ \frac{3.06 - 2.72}{3.06} \right]$$
  

$$\therefore 10 \times 3.06 = l_1 (0.34)$$
  

$$\therefore 30.6 = 0.34 l_1$$
  

$$\therefore l_1 = \frac{30.6}{0.34} = 90$$
  
Now,  $d_1 = l_1 - l_2$   

$$\therefore 10 = 90 - l_2$$
  

$$\therefore l_2 = 90 - 10 \qquad \therefore l_2 = 80$$

Hence, the number of persons living at age 1 is  $l_1 = 90$  and the number of persons living at age 2 is  $l_2 = 80$ .

(4 MARKS)

#### ANSWER 2(B)

Jobs	Department		
	Α	В	
I	8	8	
II	6	3	
	5	4	

**Step I :** Min.  $(M_{i1}, M_{i2}) = 3$ , which corresponds to department B. Therefore, job II is processed in the last

The problem now reduces to two jobs I and III.

Here, Min.  $(M_{i1}, M_{i2}) = 4$ , which corresponds to the department B.

Therefore, job III is processes in the last next to job III and then job I is processed. Thus, the optimal sequence of jobs is obtained as follows :

I III II
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Minimum elapsed time can be computed as follows :

Job	Department A		Depar	Idle time	
	Time in	Time out	Time in	Time out	for B
I	0	8	8	16	8
III	8	13	16	20	0
=	13	19	20	23	0
		8			

From the above table,

Minimum total elapsed time to finish all three jobs T = 23 days.

Idle time for the department A

= T - {sum of the processing time to finish all jobs in A}

= 23 – 19

= 4 days.

Idle time for the department B = 8 days.

## ANSWER 3(A)

Since age – group wise information is given, we have

 $\mathsf{CDR} = \frac{\sum D_i}{\sum P_i} \times 1000$ 

Hence,  $12.0 = \frac{\sum D_i}{\sum P_i} \times 1000$ 

Thus,  $12.0 = \frac{80+x}{7500} \times 1000$ 

i.e. 10(80 + x) = 900

i.e. 800 + 10x = 900

$$\Rightarrow$$
 10x = 100

∴ x = 10

(2 MARKS)

## ANSWER 3(B)

**Step 1** : We replace – by very large number say  $\infty$ . No. of antibiotic product and No. of capsulation machines are not same. It is balanced by introduction of dummy machine C<sub>5</sub> with zero cost. We get

Capsulation	Antibiotic Products						
Machines	Α	В	С	D	E		
C <sub>1</sub>	27	18	8 S	20	21		
C <sub>2</sub>	31	24	21	12	17		
C <sub>3</sub>	20	17	20	$\infty$	16		
C <sub>4</sub>	21	28	20	16	27		
C <sub>5</sub>	0	0	0	0	0		

Step 2 :	Subtract the	smallest elemen	t in each row fror	m every element in	that row. We get
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Capsulation	Antibiotic Products						
Machines	Α	В	С	D	Ε		
C <sub>1</sub>	9	10	8	2	3		
C <sub>2</sub>	19	12	9	φ	5		
C₃	-4	1	4		-0		
$C_4$	5	12	4	φ	11		
C <sub>5</sub>	-0	0	0	•	-0		

(4 MARKS)

**Step 3**: Since, there is one zero in each column, the reduced matrix obtained in Step 2 remains unchanged.

- **Step 4:** Since, the number of straight lines covering all zeros is not equal to the number of rows / columns, the optimal solution has not reached.
- **Step 5 :** Subtract the smallest element 4 among the uncovered elements from each of the uncovered elements and add it to the elements at the intersection of two lines. We get,

Capsulation	Antibiotic Products				
Machines	Α	В	С	D	E
C1	9	φ	<del>6</del>	6	8
C <sub>2</sub>	15	8	5	¢	1
C <sub>3</sub>	4	1	4	æ	þ
$C_4$	1	8	þ	þ	
C <sub>5</sub>	0	6	6	4	6

**Step 6:** Since, the number of straight lines covering all zeros is equal to the number of rows / columns, the optimum solution has reached. The optimal assignment can be made as follows :

Capsulation		Antibio	otic Pro	ducts	
Machines	Α	В	С	D	Ε
<b>C</b> <sub>1</sub>	9	0	8	6	3
C <sub>2</sub>	15	8	5	0	1
C <sub>3</sub>	4	1	4	$\infty$	0
$C_4$	1	8	0	X	7
C <sub>5</sub>	0	X	X	4	X

Hence, optimal assignment schedule is obtained as follows :

Antibiotic Products	Capsulation machines	Cost (Rs.)
А	C <sub>5</sub>	0
В	C <sub>1</sub>	18
С	$C_4$	20
D	C <sub>2</sub>	12
E	C <sub>3</sub>	16

Total (minimum) cost = 18 + 20 + 12 + 16 = Rs. 66.

**[Optimum assignment schedule :**  $M_1 \rightarrow B$ ,  $M_2 \rightarrow D$ ,  $M_3 \rightarrow E$ ,  $M_4 \rightarrow C$ ,  $M_5 \rightarrow A$ .

Minimum time required = 66.]

(4 MARKS)

#### ANSWER 3(C)

Given :  $\Sigma D$  = 812,  $\Sigma P$  = 80000

Now, CDR = 
$$\frac{\sum D}{\sum P} \times 1000$$
  
=  $\frac{812}{80000} \times 1000$   
=  $\frac{812}{80} = 10.15$ 

∴ CDR = 10.15

Hence, CDR = 10.15 per thousand.

(2 MARKS)

## ANSWER 4(A)

x	q <sub>x</sub>	$P_x = 1 - q_x$	$l_x$	$L_{x} = \frac{l_{x} + l_{x+1}}{2}$	$\mathbf{T}_{\mathbf{x}} = \mathbf{L}_{\mathbf{x}} + \mathbf{T}_{\mathbf{x}+1}$
0	0.3	0.7	1000	850	2757.36
1	0.1	0.9	700	665	1907.36
2	0.2	0.8	630	567	1242.36
3	0.4	0.6	504	403.2	675.36
4	0.6	0.4	302.4	211.68	272.16
5	1.0	0.0	120.96	60.48	60.48

Given that  $l_0 = 1000$  and the entire column  $q_x$  is specifically mortality rate. To answer question (i) we need to find  $T_x$ . We construct the life table as follows :

Note that,

 $l_{x+1} = l_x - d_x = l_x - q_x l_x = l_x (1 - q_x) = l_x P_x$ 

to find L<sub>5</sub>, we take  $I_6 = 0$ 

to compute  $T_x$ , let  $T_5 = L_5 + L_6 + .... = L_5$ 

Then using the relation,  $T_x = L_x + T_{x+1}'$ 

We compute T<sub>x</sub>

(i) The total population of birds

 $T_0 = 2757$ 

- (ii) The estimated population between the ages 1 and 4 is  $= T_1 - T_5$ = 1907.36 - 60.48 = 1846.88 = 1847
- (iii) With intake of 1000  $(l_0)$  we get total population 2757.36, thus to get total population to be 4000 we need to take proportionately more intake.

old intake	New intake				
Total old population	$=$ $\frac{1}{Total new population}$				
$\frac{1000}{2757.36} =$	New intake 4000				
∴ New intake = 1450.662952					
Thus, Additional intake = ( <i>New intake</i> - <i>old intake</i> ) = 1450.66 - 1000 = 450.66 $\approx$ 451 Thus 451 new born birds need to be added for required purpose.					

(4 MARKS)

## ANSWER 4(B)

**Step : 1 :** "Express the problem into  $n \times n$ , where n = 4 assignment matrix". This step has already been done.

**Step 2 : "**Select the smallest element in each row and subtract it from every element in its row so as to obtain atleast one zero in each row".

25, 21, 19 and 34 are the smallest elements in the first, second, third and fourth row respectively. Subtraction of the smallest element from every element in its row gives rise to a new assignment matrix.

	Table					
	Machines					
Jobs	Р	Q	R	S		
	F	Processing cost (Rs.)				
Α	6	6 0 8 4				
В	4	3	2	0		
C	0	2	4	5		
D	4	2	0	6		

**Step 3: "**Select the smallest element in each column of the reduced matrix obtained in step 2 and subtract it from every element in that column".

Since each column in table contains one Zero, subtraction of minimum element in each column from every element in that column will not make further change in the assignment matrix of table.

Step 4 : Cover all zeros by minimum number of straight lines (horizontal and vertical lines only)

Minimum four lines are required to cover all zeros



As the minimum number of straight lines required to connect all zeros in the assignment matrix equals n (number of rows/ columns), optimal solution has reached.

**Step 5 : (i)** Examine the rows one by one starting with the first row until a row with exactly one zero is found. Mark the zero by enclosing it in square ( $\Box$ ), indicating assignment of the job. Cross all the zeros in the same column as they cannot be used to make other assignment.

(ii) "Examine next the columns for any column with exactly one zero and mark each as above, crossing the remaining zeros in that row."

This step is shown in the Table

	Machines				
Jobs	Р	Q	R	S	
	Р	rocessing	cost (R	s.)	
A	6	0	8	4	
В	4	3	2	0	
C	0	2	4	5	
D	4	2	0	6	

It is observed that all the zeros are assigned and each row and each column contains exactly one assignment. Hence the optimal (Minimum) assignment schedule is displayed in table. Note that choices of horizontal and vertical lines as shown in table need not be same. However the final solution will be the same.

₹s.)

Hence total (Minimum cost is = 25 + 21 + 19 + 34 = Rs. 99.

## ANSWER 5(A)

$$\mathsf{CDR} = \frac{\sum D_i}{\sum P_i} \times 1000$$

For population A :

 $\Sigma D_i = 170 + 115 + 490 + 630 = 1405$ 

 $\Sigma P_i = 13 + 20 + 52 + 22 = 107$  (in thousands)

 $\therefore\,$  CDR for population A denoted by  $\text{CDR}_{\text{A}}$  is

$$CDR_A = \frac{\sum D_i}{\sum P_i} \times 1000$$

$$=\frac{1405}{107000} \times 1000$$

=13.13 per thousand.

For population B :

 $\Sigma D_i = 510 + 130 + 570 + 680 = 1890$ 

 $\Sigma P_i = 15 + 35 + 54 + 23 = 127$  (in thousands)

 $\therefore$  CDR for population B denoted by  $\mathsf{CDR}_{\mathsf{B}}$  is,

$$\mathsf{CDR}_{\mathsf{B}} = \frac{\sum D_i}{\sum P_i} \times 100$$

$$=\frac{1890}{127000} \times 1000$$

= 14.88 per thousand.

Observe that population A is more healthy than population B as  $CDR_A < CDR_B$ .

#### (3 MARKS)

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(4 MARKS)

## ANSWER 5(B)

**Step 1**: "Express the problem into  $n \times n$ , where n = 4 assignment matrix." This step has already been done.

Step 2 : "Subtract the smallest element in each row from every element in its row".

This	This step is shown in Table				
			Table	-1	
	Jobs		Mach	ines	
		M1	M <sub>2</sub>	M₃	$M_4$
	$J_1$	0	10	20	25
	J <sub>2</sub>	0	8	16	18
	$J_3$	0	8	16	18
	$J_4$	0	6	12	20

**Step 3 :** "Subtract the smallest element in each column of the assignment matrix obtained in step 2 from every element in that column".

		Table	- 2	
Jobs		Mac	hines	
	M1	M <sub>3</sub>	$M_4$	
$J_1$	0	4	8	7
J <sub>2</sub>	0	2	4	0
J <sub>3</sub>	0	2	4	0
$J_4$	0	0	0	2

Step 4: "Cover all zeros with minimum number of straight lines".

		lable	-3	
Jobs		Mac	hines	
	M1	M <sub>2</sub>	M <sub>3</sub>	M <sub>4</sub>
$J_1$	φ	4	8	7
J <sub>2</sub>	φ	2	4	Ø
$J_3$	φ	2	4	þ
$J_4$	<del>0</del>	0	0	2

As the number of lines required to cover all zeros is less than the number of rows / columns, optimal solution has still not reached.

**Step 5**: "Select the smallest element not covered by the lines from table 3, subtract it from each uncovered element and add it to elements which are at the intersection of the lines". This step is shown in Table 4.

		Tabl	e 4	
Jobs	Machines			
	M <sub>1</sub> M <sub>2</sub> M <sub>3</sub> M			
J <sub>1</sub>	0	2	6	7
J <sub>2</sub>	0	0	2	0
J <sub>3</sub>	0	0	2	0
$J_4$	2	0	0	4

Step 6 : "Cover all zeros by minimum number of straight lines".

Table 5
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Jobs		Mac	hines	
	<b>M</b> <sub>1</sub>	M <sub>2</sub>	M <sub>3</sub>	M <sub>4</sub>
$J_1$	φ	2	6	7
J <sub>2</sub>	φ	Ø	2	þ
J <sub>3</sub>	φ	Ø	2	þ
$J_4$	2	0	0	4

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Since the number of lines covering all zeros is equal to the number of rows / columns, hence the optimal solution has reached. Optimal solution can be made as follows :

		Table	e 6		
Jobs	Machines				
	M1	M <sub>2</sub>	M₃	$M_4$	
$J_1$	0	2	6	7	
J <sub>2</sub>	×	0	2	0	
J <sub>3</sub>	×	0	2	0	
$J_4$	2	X	0	4	

Note that, from table 6, row 1 has only one zero hence job  $J_1$  is assigned to machine  $M_1$  and remaining zeros in the column have been crossed out.

Similarly column  $M_3$  has only zero and as such job  $J_4$  is assigned to machine  $M_3$  and remaining zeros in row  $J_4$  have been crossed out.

For remaining zeros, the assignment can be made in two ways :

- (i) To make assignment considering the zero element in row  $J_2$  and Column  $M_2$ . If it is assigned, then the additional zero in column  $M_2$  has to cross out. This will automatically assign job  $J_3$  to Machine  $M_4$  OR
  - (ii) To make an assignment considering zero element in row  $J_2$  and column  $M_4$  and cross out other zero appearing in the column  $M_4$ . This will automatically assign job  $J_3$  to Machine  $M_2$ .

Jobs	Machines			
	M1	M <sub>2</sub>	M <sub>3</sub>	M <sub>4</sub>
J <sub>1</sub>	0	2	6	7
J <sub>2</sub>	X	0	2	×
$J_3$	X	×	2	0
$J_4$	2	×	0	4

		Т	able 8	
Jobs		Mac	hines	
	M1	M <sub>2</sub>	M₃	M <sub>4</sub>
J <sub>1</sub>	0	2	6	7
J <sub>2</sub>	×	X	2	0
J <sub>3</sub>	øX	0	2	X
$J_4$	2	×	0	4

Hence the optimal assignment and optimal cost can be computed as follows :

Jobs	Machines	Cost
$J_1$	M <sub>1</sub>	40
J <sub>2</sub>	M <sub>2</sub>	38
$J_3$	$M_4$	43
$J_4$	M <sub>3</sub>	51

Jobs	Machines	Cost
$J_1$	M <sub>1</sub>	40
J <sub>2</sub>	M <sub>4</sub>	48
$J_3$	M <sub>2</sub>	33
J <sub>4</sub>	M <sub>3</sub>	51

Total (minimum) cost

Total (Minimum) cost

= 40 + 38 + 43 + 51 = 40 + 48 + 33 + 51= 172 = 172

In both the cases, the optimal cost is Rs. 172.

(5 MARKS)

## ANSWER 5(C)

Given : $e_{40}^0 = 31$ , $I_{40} = 550$ ,		
We know that,	$e_x^o = \frac{T_x}{l_x}$	
$\therefore e_{40}^{o} = \frac{T_{40}}{l_{40}} \qquad \therefore$	$31 = \frac{T_{40}}{550}$	: $550 \times 31 = T_{40}$
∴ T <sub>40</sub> = 17050	Hence, T <sub>40</sub> = 170	050.

(2 MARKS)