# Jio TEST SERIES Evaluate Learn Succeed 

## SUGGESTED SOLUTION

SYJC<br>SUBJECT- MATHS \& STATS<br>Test Code - SyJ 6097<br>BRANCH - () (Date :)

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## ANSWER 1(A)

## Given: $I_{2}=80, d_{2}=5, \mathrm{q}_{3}=0.20, I_{3}=?, I_{4}=$ ?

We have, $\mathrm{d}_{\mathrm{x}}=I_{\mathrm{x}}-I_{\mathrm{x}+1}$
$\therefore \mathrm{d}_{2}=I_{2}-I_{3}$
$\therefore 5=80-I_{3}$
$\therefore I_{3}=80-5=75$

We have, $\mathrm{q}_{\mathrm{x}}=\frac{d_{x}}{l_{x}}$

$$
\begin{array}{ll}
\therefore & \mathrm{q}_{3}=\frac{d_{3}}{l_{3}} \\
\therefore & 0.20=\frac{d_{3}}{75} \\
\therefore & 0.20 \times 75=\mathrm{d}_{3} \\
\therefore & \mathrm{~d}_{3}=15
\end{array}
$$

Now, $d_{3}=I_{3}-I_{4}$

$$
\begin{array}{ll}
\therefore & 15=75-I_{4} \\
\therefore & I_{4}=75-15 \\
\therefore & I_{4}=60
\end{array}
$$

Hence, $I_{3}=75$ and $I_{4}=60$
(3 MARKS)

## ANSWER 1(B)

Using the optimal sequence algorithm, we can obtain the following optimal sequence.

| E | B | D | C | A |
| :--- | :--- | :--- | :--- | :--- |

Next, we find total elapsed time as follows :

| Job | Machine $\mathbf{M}_{\mathbf{1}}$ |  | Machine $\mathbf{M}_{\mathbf{2}}$ |  | Idle time for <br> $\mathbf{M}_{\mathbf{2}}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | Time in | Time out | Time in | Time Out |  |
| E | 0 | 6 | 6 | 17 | 02 |
| B | 6 | 19 | 19 | 34 | 00 |
| D | 19 | 31 | 34 | 40 | 00 |
| C | 31 | 38 | 40 | 45 | 00 |
| A | 38 | 42 | 45 | 48 | 00 |
|  |  |  |  | Total | 08 |

Minimum elapsed (Processing) time $\mathrm{T}=48 \mathrm{hrs}$.

Idle time for machine $\mathrm{M}_{1}$
$=T-\left(\right.$ sum of the processing time of the five jobs on $\left.M_{1}\right)$
$=48-42=06$ hrs.
Idle time for machine $M_{2}$
$=T-\left(\right.$ sum of the processing time of the five jobs on $\left.M_{2}\right)$
$=48-40=08$ hrs.

## ANSWER 2(A)

Let $l_{\mathrm{x}}=$ number of persons living at age x
We know that,
$\mathrm{d}_{\mathrm{x}}=l_{x}-l_{\mathrm{x}+1}$
$p_{\mathrm{x}}=\frac{l_{x+1}}{l_{x}}$ and also $p_{\mathrm{x}}=\frac{e_{x}}{1+e_{x+1}}$
$\therefore \frac{l_{x+1}}{l_{x}}=p_{\mathrm{x}}=\frac{e_{x}}{1+e_{x+1}}$
$\therefore l_{\mathrm{x}+1}=\frac{e_{x} \cdot l_{x}}{1+e_{x+1}}$
Using this result in (1) we get
$\mathrm{d}_{\mathrm{x}}=l_{\mathrm{x}}=\frac{e_{x} \cdot l_{x}}{1+e_{x+1}}$
$\therefore \mathrm{d}_{\mathrm{x}}=l_{\mathrm{x}}\left[1-\frac{e_{x}}{1+e_{x+1}}\right]$
Given : $\mathrm{d}_{1}=10, e_{1}^{0}=3.22, e_{2}^{0}=2.56$
We have $e_{x}^{0}=\mathrm{e}_{\mathrm{x}}+\frac{1}{2}$
$\therefore \mathrm{e}_{\mathrm{x}}=e_{x}^{0}-\frac{1}{2}=e_{x}^{0}-0.5$
$e_{1}^{0}=3.22$ and $e_{2}^{0}=2.56$
Using the result (3) we get
$e_{1}=3.22-0.5=2.72$
$e_{2}=2.56-0.5=2.06$

From (2),
$\mathrm{d}_{1}=l_{1}\left[1-\frac{e_{1}}{1+e_{2}}\right]$

Putting $d_{1}=10, e_{1}=2.72$ and $e_{2}=2.06$, we get
$10=l_{1}\left[1-\frac{2.72}{1+2.06}\right]$
$\therefore 10=l_{1}\left[1-\frac{2.72}{3.06}\right]$
$\therefore 10=l_{1}\left[\frac{3.06-2.72}{3.06}\right]$
$\therefore 10 \times 3.06=l_{1}(0.34)$
$\therefore 30.6=0.34 l_{1}$
$\therefore l_{1}=\frac{30.6}{0.34}=90$

Now, $\mathrm{d}_{1}=l_{1}-l_{2}$
$\therefore 10=90-l_{2}$
$\therefore l_{2}=90-10 \quad \therefore l_{2}=80$
$\therefore l_{1}=90$ and $l_{2}=80$

Hence, the number of persons living at age 1 is $l_{1}=90$ and the number of persons living at age 2 is $l_{2}=80$.
(4 MARKS)

## ANSWER 2(B)

| Jobs | Department |  |
| :---: | :---: | ---: |
|  | A | B |
| I | 8 | 8 |
| II | 6 | 3 |
| III | 5 | 4 |

Step I: Min. $\left(\mathrm{M}_{\mathrm{i} 1}, \mathrm{M}_{\mathrm{i} 2}\right)=3$, which corresponds to department B.
Therefore, job II is processed in the last


The problem now reduces to two jobs I and III.
Here, Min. $\left(M_{i 1}, M_{i 2}\right)=4$, which corresponds to the department $B$.

Therefore, job III is processes in the last next to job III and then job I is processed. Thus, the optimal sequence of jobs is obtained as follows :

| I | III | II |
| :--- | :--- | :--- |

Minimum elapsed time can be computed as follows :

| Job | Department A |  | Department B |  | Idle time |
| :---: | :---: | :---: | :---: | :---: | :---: |
| for B |  |  |  |  |  |

From the above table,

Minimum total elapsed time to finish all three jobs T=23 days.

Idle time for the department $A$
$=\mathrm{T}-\{$ sum of the processing time to finish all jobs in A$\}$
$=23-19$
$=4$ days.
Idle time for the department $B=8$ days.
(4 MARKS)

## ANSWER 3(A)

Since age - group wise information is given, we have

$$
\begin{aligned}
& \quad C D R=\frac{\sum D_{i}}{\sum P_{i}} \times 1000 \\
& \text { Hence, } \quad 12.0=\frac{\sum D_{i}}{\sum P_{i}} \times 1000 \\
& \text { Thus, } \quad 12.0=\frac{80+x}{7500} \times 1000 \\
& \text { i.e. } \quad 10(80+\mathrm{x})=900 \\
& \text { i.e. } \quad 800+10 \mathrm{x}=900 \\
& \Rightarrow \quad 10 \mathrm{x}=100 \\
& \\
& \therefore \quad \mathrm{x}=10
\end{aligned}
$$

(2 MARKS)

## ANSWER 3(B)

Step 1 : We replace - by very large number say $\infty$. No. of antibiotic product and No. of capsulation machines are not same. It is balanced by introduction of dummy machine $C_{5}$ with zero cost. We get

| Capsulation | Antibiotic Products |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Machines | A | B | C | $\mathbf{D}$ | E |
| $\mathrm{C}_{1}$ | 27 | 18 | $\infty$ | 20 | 21 |
| $\mathrm{C}_{2}$ | 31 | 24 | 21 | 12 | 17 |
| $\mathrm{C}_{3}$ | 20 | 17 | 20 | $\infty$ | 16 |
| $\mathrm{C}_{4}$ | 21 | 28 | 20 | 16 | 27 |
| $\mathrm{C}_{5}$ | 0 | 0 | 0 | 0 | 0 |

Step 2 : Subtract the smallest element in each row from every element in that row. We get

| Capsulation | Antibiotic Products |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Machines | A | B | C | D | E |
| $\mathrm{C}_{1}$ | 9 | 10 | $\infty$ | 2 | 3 |
| $\mathrm{C}_{2}$ | 19 | 12 | 9 | $\oint$ | 5 |
| $\mathrm{C}_{3}$ | 4 | 1 | 4 |  | 0 |
| $\mathrm{C}_{4}$ | 5 | 12 | 4 | $\oint$ | 11 |
| $\mathrm{C}_{5}$ | 0 | 0 | 0 | $\oint$ | 0 |

Step 4: Since, the number of straight lines covering all zeros is not equal to the number of rows / columns, the optimal solution has not reached.

Step 5 : Subtract the smallest element 4 among the uncovered elements from each of the uncovered elements and add it to the elements at the intersection of two lines. We get,

| Capsulation | Antibiotic Products |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Machines | A | B | C | D | E |  |
| $\mathrm{C}_{1}$ | 9 | 0 | 0 | $\$$ | $\beta$ |  |
| $\mathrm{C}_{2}$ | 15 | 8 | 5 | 0 | 1 |  |
| $\mathrm{C}_{3}$ | 4 | 1 | 4 | $\phi$ | 0 |  |
| $\mathrm{C}_{4}$ | 1 | 8 | 0 | 0 | $\$$ |  |
| $\mathrm{C}_{5}$ | 0 | 0 | 0 | 4 | 0 |  |

Step 6: Since, the number of straight lines covering all zeros is equal to the number of rows / columns, the optimum solution has reached. The optimal assignment can be made as follows :

| Capsulation | Antibiotic Products |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Machines | A | B | C | D | E |
| $\mathrm{C}_{1}$ | 9 | 0 | $\infty$ | 6 | 3 |
| $\mathrm{C}_{2}$ | 15 | 8 | 5 | 0 | 1 |
| $\mathrm{C}_{3}$ | 4 | 1 | 4 | $\infty$ | 0 |
| $\mathrm{C}_{4}$ | 1 | 8 | 0 | $\npreceq$ | 7 |
| $\mathrm{C}_{5}$ | 0 | $\not 2$ | $\not \not 2$ | 4 | $\npreceq$ |

Hence, optimal assignment schedule is obtained as follows :

| Antibiotic Products | Capsulation machines | Cost (Rs.) |
| :---: | :---: | :---: |
| A | $\mathrm{C}_{5}$ | 0 |
| B | $\mathrm{C}_{1}$ | 18 |
| C | $\mathrm{C}_{4}$ | 20 |
| D | $\mathrm{C}_{2}$ | 12 |
| E | $\mathrm{C}_{3}$ | 16 |

Total ( minimum) cost $=18+20+12+16=$ Rs. 66.
[Optimum assignment schedule : $M_{1} \rightarrow B, M_{2} \rightarrow D, M_{3} \rightarrow E, M_{4} \rightarrow C, M_{5} \rightarrow A$.
Minimum time required $=66$.]
(4 MARKS)

## ANSWER 3(C)

Given : $\Sigma \mathrm{D}=812, \Sigma \mathrm{P}=80000$
Now, CDR $=\frac{\sum D}{\sum P} \times 1000$

$$
\begin{aligned}
& =\frac{812}{80000} \times 1000 \\
& =\frac{812}{80}=10.15
\end{aligned}
$$

$\therefore C D R=10.15$
Hence, CDR = 10.15 per thousand.

## ANSWER 4(A)

Given that $l_{0}=1000$ and the entire column $\mathrm{q}_{\mathrm{x}}$ is specifically mortality rate. To answer question (i) we need to find $T_{x}$. We construct the life table as follows :

| $\mathbf{x}$ | $\mathbf{q}_{\mathbf{x}}$ | $\mathbf{P}_{\mathbf{x}}=\mathbf{1}-\mathbf{q}_{\mathbf{x}}$ | $\boldsymbol{l}_{\boldsymbol{x}}$ | $\mathbf{L}_{\mathbf{x}}=\frac{\boldsymbol{l}_{\boldsymbol{x}}+\boldsymbol{l}_{\mathbf{x + 1}}}{\mathbf{2}}$ | $\mathbf{T}_{\mathrm{x}}=\mathrm{L}_{\mathbf{x}}+\mathbf{T}_{\mathrm{x}+\mathbf{1}}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0.3 | 0.7 | 1000 | 850 | 2757.36 |
| 1 | 0.1 | 0.9 | 700 | 665 | 1907.36 |
| 2 | 0.2 | 0.8 | 630 | 567 | 1242.36 |
| 3 | 0.4 | 0.6 | 504 | 403.2 | 675.36 |
| 4 | 0.6 | 0.4 | 302.4 | 211.68 | 272.16 |
| 5 | 1.0 | 0.0 | 120.96 | 60.48 | 60.48 |

Note that,
$l_{\mathrm{x}+1}=l_{\mathrm{x}}-\mathrm{d}_{\mathrm{x}}=l_{\mathrm{x}}-\mathrm{q}_{\mathrm{x}} l_{\mathrm{x}}=l_{\mathrm{x}}\left(1-\mathrm{q}_{\mathrm{x}}\right)=l_{\mathrm{x}} \mathrm{P}_{\mathrm{x}}$
to find $L_{5}$, we take $I_{6}=0$
to compute $\mathrm{T}_{\mathrm{x}}$, let $\mathrm{T}_{5}=\mathrm{L}_{5}+\mathrm{L}_{6}+\ldots . .=\mathrm{L}_{5}$

Then using the relation, $\mathrm{T}_{\mathrm{x}}=\mathrm{L}_{\mathrm{x}}+\mathrm{T}_{\mathrm{x}+1}{ }^{\prime}$

We compute $T_{x}$
(i) The total population of birds
$\mathrm{T}_{0}=2757$
(ii) The estimated population between the ages 1 and 4 is
$=T_{1}-T_{5}$
= $1907.36-60.48$
$=1846.88=1847$
(iii) With intake of $1000\left(l_{0}\right)$ we get total population 2757.36 , thus to get total population to be 4000 we need to take proportionately more intake.

$$
\begin{aligned}
\frac{\text { old intake }}{\text { Total old population }} & =\frac{\text { New intake }}{\text { Total new population }} \\
\frac{1000}{2757.36} & =\frac{\text { New intake }}{4000}
\end{aligned}
$$

$\therefore$ New intake $=1450.662952$

Thus, Additional intake $=($ New intake - old intake $)$

$$
\begin{aligned}
& =1450.66-1000 \\
& =450.66 \approx 451
\end{aligned}
$$

Thus 451 new born birds need to be added for required purpose.
(4 MARKS)

Step : 1 : "Express the problem into $n \times n$, where $n=4$ assignment matrix". This step has already been done.
Step 2 : "Select the smallest element in each row and subtract it from every element in its row so as to obtain atleast one zero in each row".

25, 21, 19 and 34 are the smallest elements in the first, second, third and fourth row respectively. Subtraction of the smallest element from every element in its row gives rise to a new assignment matrix.

Table

|  | Machines |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| Jobs | $\mathbf{P}$ | $\mathbf{Q}$ | $\mathbf{R}$ | $\mathbf{S}$ |
|  | Processing cost (Rs.) |  |  |  |
| A | 6 | 0 | 8 | 4 |
| B | 4 | 3 | 2 | 0 |
| C | 0 | 2 | 4 | 5 |
| D | 4 | 2 | 0 | 6 |

Step 3: "Select the smallest element in each column of the reduced matrix obtained in step 2 and subtract it from every element in that column".

Since each column in table contains one Zero, subtraction of minimum element in each column from every element in that column will not make further change in the assignment matrix of table.

Step 4 : Cover all zeros by minimum number of straight lines (horizontal and vertical lines only)
Minimum four lines are required to cover all zeros


As the minimum number of straight lines required to connect all zeros in the assignment matrix equals $n$ (number of rows/ columns), optimal solution has reached.
Step 5 : (i) Examine the rows one by one starting with the first row until a row with exactly one zero is found. Mark the zero by enclosing it in square ( $\square$ ), indicating assignment of the job. Cross all the zeros in the same column as they cannot be used to make other assignment.
(ii) "Examine next the columns for any column with exactly one zero and mark each as above, crossing the remaining zeros in that row."
This step is shown in the Table
Table

|  | Machines |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| Jobs | $\mathbf{P}$ | $\mathbf{Q}$ | $\mathbf{R}$ | $\mathbf{S}$ |
|  | Processing cost (Rs.) |  |  |  |
| A | 6 | 0 | 8 | 4 |
| B | 4 | 3 | 2 | 0 |
| C | $\boxed{0}$ | 2 | 4 | 5 |
| D | 4 | 2 | $\boxed{0}$ | 6 |

It is observed that all the zeros are assigned and each row and each column contains exactly one assignment. Hence the optimal (Minimum) assignment schedule is displayed in table. Note that choices of horizontal and vertical lines as shown in table need not be same. However the final solution will be the same.

Table

| Job | Machines | Cost (Rs.) |
| :---: | :---: | :---: |
| A | Q | 25 |
| B | S | 21 |
| C | P | 19 |
| D | R | 34 |

Hence total (Minimum cost is =

$$
25+21+19+34=\text { Rs. } 99 .
$$

## ANSWER 5(A)

$\mathrm{CDR}=\frac{\sum D_{i}}{\sum P_{i}} \times 1000$
For population A :
$\Sigma \mathrm{D}_{i}=170+115+490+630=1405$
$\Sigma \mathrm{P}_{\mathrm{i}}=13+20+52+22=107$ (in thousands)
$\therefore$ CDR for population $A$ denoted by $\operatorname{CDR}_{\mathrm{A}}$ is

$$
\begin{aligned}
\mathrm{CDR}_{\mathrm{A}}= & \frac{\sum D_{i}}{\sum P_{i}} \times 1000 \\
& =\frac{1405}{107000} \times 1000 \\
& =13.13 \text { per thousand. } .
\end{aligned}
$$

## For population B :

$\Sigma D_{i}=510+130+570+680=1890$
$\Sigma \mathrm{P}_{\mathrm{i}}=15+35+54+23=127$ (in thousands)
$\therefore$ CDR for population B denoted by $\mathrm{CDR}_{\mathrm{B}}$ is,

$$
\begin{aligned}
\mathrm{CDR}_{\mathrm{B}}= & \frac{\sum D_{i}}{\sum P_{i}} \times 100 \\
= & \frac{1890}{127000} \times 1000 \\
= & 14.88 \text { per thousand. }
\end{aligned}
$$

Observe that population A is more healthy than population B as $\mathrm{CDR}_{\mathrm{A}}<\mathrm{CDR}_{\mathrm{B}}$.

## ANSWER 5(B)

Step 1 : "Express the problem into $\mathrm{n} \times \mathrm{n}$, where $\mathrm{n}=4$ assignment matrix." This step has already been done.
Step 2 : "Subtract the smallest element in each row from every element in its row".

This step is shown in Table
Table - 1

| Jobs | Machines |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
|  | $\mathbf{M}_{\mathbf{1}}$ | $\mathbf{M}_{\mathbf{2}}$ | $\mathbf{M}_{\mathbf{3}}$ | $\mathbf{M}_{\mathbf{4}}$ |
| $\mathbf{J}_{\mathbf{1}}$ | 0 | 10 | 20 | 25 |
| $\mathbf{J}_{\mathbf{2}}$ | 0 | 8 | 16 | 18 |
| $\mathbf{J}_{\mathbf{3}}$ | 0 | 8 | 16 | 18 |
| $\mathbf{J}_{\mathbf{4}}$ | 0 | 6 | 12 | 20 |

Step 3 : "Subtract the smallest element in each column of the assignment matrix obtained in step 2 from every element in that column".

Table-2

| Jobs | Machines |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
|  | $\mathbf{M}_{\mathbf{1}}$ | $\mathbf{M}_{\mathbf{2}}$ | $\mathbf{M}_{\mathbf{3}}$ | $\mathbf{M}_{\mathbf{4}}$ |
| $\mathbf{J}_{\mathbf{1}}$ | 0 | 4 | 8 | 7 |
| $\mathbf{J}_{\mathbf{2}}$ | 0 | 2 | 4 | 0 |
| $\mathbf{J}_{\mathbf{3}}$ | 0 | 2 | 4 | 0 |
| $\mathbf{J}_{\mathbf{4}}$ | 0 | 0 | 0 | 2 |

Step 4 : "Cover all zeros with minimum number of straight lines".
Table-3

| Jobs | Machines |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | $\mathbf{M}_{\mathbf{1}}$ | $\mathbf{M}_{\mathbf{2}}$ | $\mathbf{M}_{\mathbf{3}}$ | $\mathbf{M}_{\mathbf{4}}$ |  |
| $\mathbf{J}_{\mathbf{1}}$ | $\oint$ | 4 | 8 |  |  |
| $\mathbf{J}_{\mathbf{2}}$ | $\oint$ | 2 | 4 | 0 |  |
| $\mathbf{J}_{\mathbf{3}}$ | $\oint$ | 2 | 4 | 0 |  |
| $\mathbf{J}_{\mathbf{4}}$ | $\oplus$ | 0 | 0 | $\mathbf{Z}$ |  |

As the number of lines required to cover all zeros is less than the number of rows / columns, optimal solution has still not reached.

Step 5 : "Select the smallest element not covered by the lines from table 3, subtract it from each uncovered element and add it to elements which are at the intersection of the lines". This step is shown in Table 4.

Table 4

| Jobs | Machines |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
|  | $\mathbf{M}_{\mathbf{1}}$ | $\mathbf{M}_{\mathbf{2}}$ | $\mathbf{M}_{\mathbf{3}}$ | $\mathbf{M}_{\mathbf{4}}$ |
| $\mathbf{J}_{\mathbf{1}}$ | 0 | 2 | 6 | 7 |
| $\mathbf{J}_{\mathbf{2}}$ | 0 | 0 | 2 | 0 |
| $\mathbf{J}_{\mathbf{3}}$ | 0 | 0 | 2 | 0 |
| $\mathbf{J}_{\mathbf{4}}$ | 2 | 0 | 0 | 4 |

Step 6 : "Cover all zeros by minimum number of straight lines".
Table 5

| Jobs | Machines |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | $\mathbf{M}_{\mathbf{1}}$ | $\mathbf{M}_{\mathbf{2}}$ | $\mathbf{M}_{\mathbf{3}}$ | $\mathbf{M}_{\mathbf{4}}$ |  |
| $\mathbf{J}_{\mathbf{1}}$ | $\oint$ | 2 | 6 | $\downarrow$ |  |
| $\mathbf{J}_{\mathbf{2}}$ | $\oint$ | $\emptyset$ | 2 | 0 |  |
| $\mathbf{J}_{\mathbf{3}}$ | $\oint$ | $\emptyset$ | 2 | $\emptyset$ |  |
| $\mathbf{J}_{\mathbf{4}}$ | $\mathbf{y}$ | $\emptyset$ | 0 | 4 |  |

Since the number of lines covering all zeros is equal to the number of rows / columns, hence the optimal solution has reached. Optimal solution can be made as follows :

Table 6

| Jobs | Machines |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
|  | $\mathbf{M}_{1}$ | $\mathrm{M}_{2}$ | $\mathrm{M}_{3}$ | M ${ }_{4}$ |
| $\mathrm{J}_{1}$ | 0 | 2 | 6 | 7 |
| $\mathrm{J}_{2}$ | $\not \not \equiv$ | 0 | 2 | 0 |
| $\mathrm{J}_{3}$ | $\nless$ | 0 | 2 | 0 |
| $\mathrm{J}_{4}$ | 2 | $x$ | 0 | 4 |

Note that, from table 6, row 1 has only one zero hence job $J_{1}$ is assigned to machine $M_{1}$ and remaining zeros in the column have been crossed out.

Similarly column $M_{3}$ has only zero and as such job $J_{4}$ is assigned to machine $M_{3}$ and remaining zeros in row $J_{4}$ have been crossed out.

For remaining zeros, the assignment can be made in two ways :
(i) To make assignment considering the zero element in row $\mathrm{J}_{2}$ and Column $\mathrm{M}_{2}$. If it is assigned, then the additional zero in column $M_{2}$ has to cross out. This will automatically assign job $J_{3}$ to Machine $M_{4}$ OR
(ii) To make an assignment considering zero element in row $\mathrm{J}_{2}$ and column $\mathrm{M}_{4}$ and cross out other zero appearing in the column $\mathrm{M}_{4}$. This will automatically assign job $\mathrm{J}_{3}$ to Machine $\mathrm{M}_{2}$.

Table 7

| Jobs | Machines |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
|  | $\mathbf{M}_{\mathbf{1}}$ | $\mathbf{M}_{\mathbf{2}}$ | $\mathbf{M}_{\mathbf{3}}$ | $\mathbf{M}_{\mathbf{4}}$ |
| $\mathbf{J}_{1}$ | $\boxed{0}$ | 2 | 6 | 7 |
| $\mathbf{J}_{\mathbf{2}}$ | $\not \nsim$ | 0 | 2 | $\nsupseteq$ |
| $\mathbf{J}_{3}$ | $\not \nsim$ | $\nsupseteq$ | 2 | $\boxed{0}$ |
| $\mathbf{J}_{4}$ | 2 | $\nsupseteq$ | $\boxed{0}$ | 4 |

Table 8

| Jobs | Machines |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
|  | $\mathrm{M}_{1}$ | $\mathbf{M}_{\mathbf{2}}$ | $\mathrm{M}_{3}$ | M ${ }_{4}$ |
| $\mathrm{J}_{1}$ | 0 | 2 | 6 | 7 |
| $\mathrm{J}_{2}$ | OX | 0 | 2 | 0 |
| $\mathrm{J}_{3}$ | OX | 0 | 2 | $\nsupseteq$ |
| $\mathrm{J}_{4}$ | 2 | $\phi$ | 0 | 4 |


| Jobs | Machines | Cost |
| :---: | :---: | :---: |
| $\mathrm{J}_{1}$ | $\mathrm{M}_{1}$ | 40 |
| $\mathrm{~J}_{2}$ | $\mathrm{M}_{2}$ | 38 |
| $\mathrm{~J}_{3}$ | $\mathrm{M}_{4}$ | 43 |
| $\mathrm{~J}_{4}$ | $\mathrm{M}_{3}$ | 51 |
|  |  |  |
|  |  |  |
|  |  |  |


| Jobs | Machines | Cost |
| :---: | :---: | :---: |
| $\mathrm{J}_{1}$ | $\mathrm{M}_{1}$ | 40 |
| $\mathrm{~J}_{2}$ | $\mathrm{M}_{4}$ | 48 |
| $\mathrm{~J}_{3}$ | $\mathrm{M}_{2}$ | 33 |
| $\mathrm{~J}_{4}$ | $\mathrm{M}_{3}$ | 51 |
|  |  |  |
|  |  |  |
|  |  |  |

Total (minimum) cost
$=40+38+43+51$
= 172
Total (Minimum) cost
$=40+48+33+51$
$=172$

In both the cases, the optimal cost is Rs. 172.

## ANSWER 5(C)

Given : $e_{40}^{0}=31, I_{40}=550$,

$$
\begin{array}{ll}
\text { We know that, } & e_{x}^{o}=\frac{T_{x}}{l_{x}} \\
\therefore & e_{40}^{o}=\frac{T_{40}}{l_{40}} \quad \therefore \\
\therefore \mathrm{~T}_{40}=17050 & \\
\begin{array}{ll} 
& \\
& \text { Hence, } \mathrm{T}_{40}=17050 .
\end{array}
\end{array}
$$

